

⇒ Levels of Programming Languages:-

→ Language Categories:-

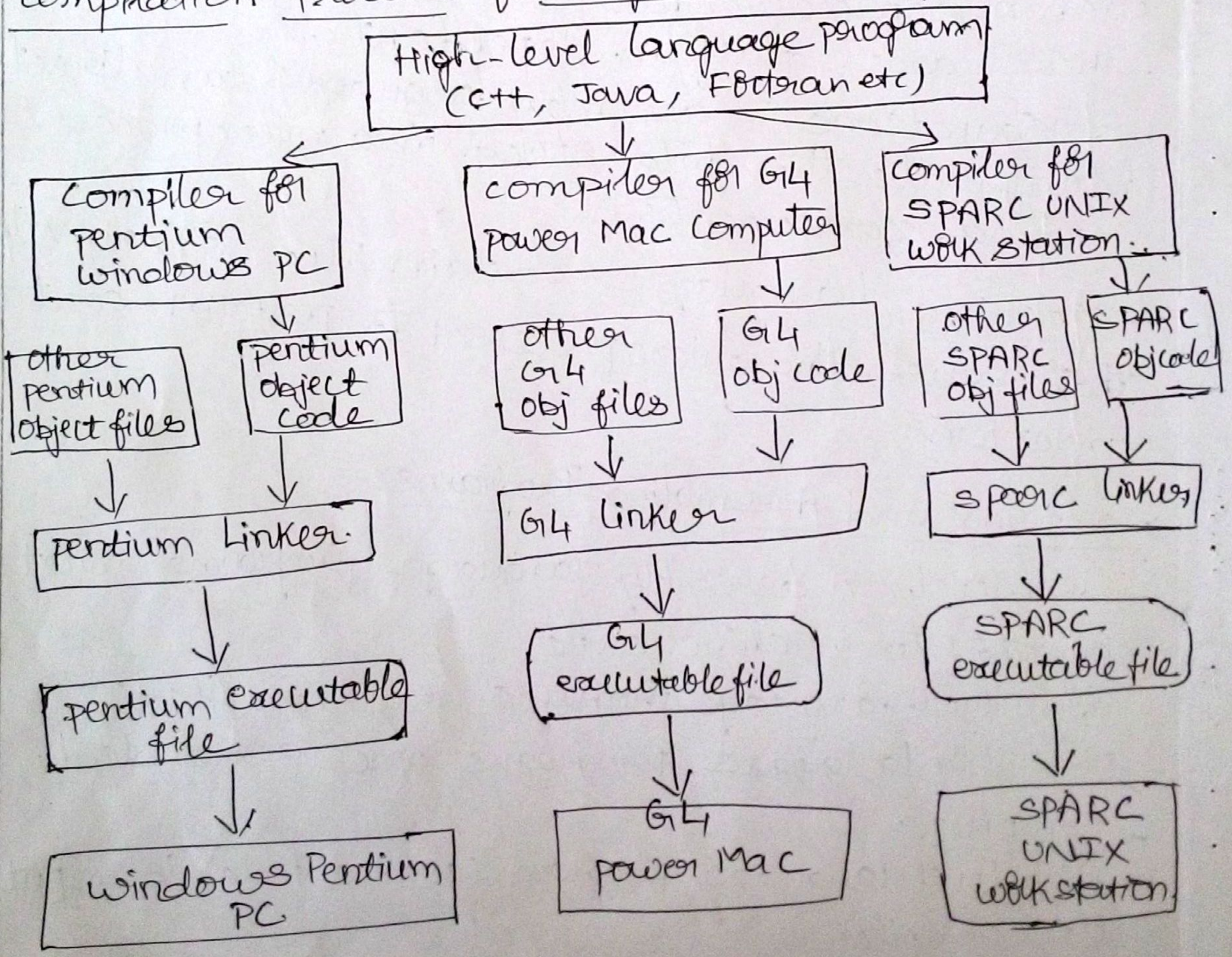
- Computer programming languages are divided into three categories.
 - High-level languages - C++, Java
 - Assembly Languages
 - Machine Languages / Low-level languages
- Languages, with the highest level of abstraction are collectively referred to as high-level languages.
- High-level languages are platform independent.
- Assembly languages are at a much ^{lower} level of abstraction. Each microprocessor has its own assembly language.
- These are platform-dependent.
- Instructions in assembly languages can directly manipulate the data stored in a microprocessor's internal components.
- Machine Languages contain the binary values that cause the microprocessor to perform certain operations.

→ Compiling and Assembling Programs:-

- High-level & assembly language programs must be converted to machine code.
- High-level language programs are compiled; assembly language programs are assembled.
- Compiler:-
 - High-level language program (source code) is input to the compiler
 - checks every statement of source code for any syntax errors; if no errors generate an object code.

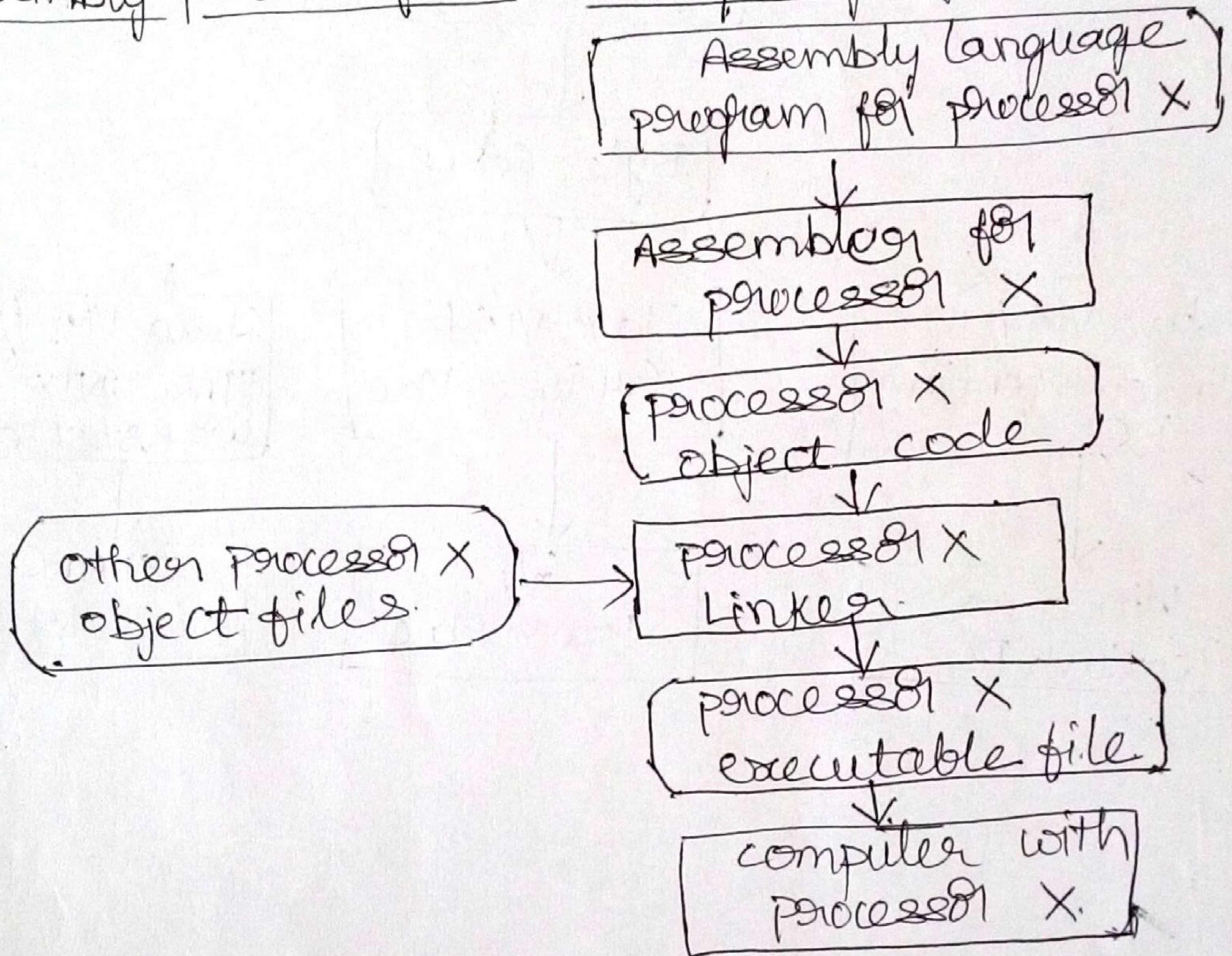
- A linker combines ^{object code} with any other ~~ob~~ required object code, to get executable file.
- A loader copies the executable file into memory. Microprocessor then runs the machine code contained in that file.
- As high-level languages are platform-independent, the same high-level source code can be compiled to run on different microprocessors and OS, & computing platforms.

Compilation Process for high-level programs:-



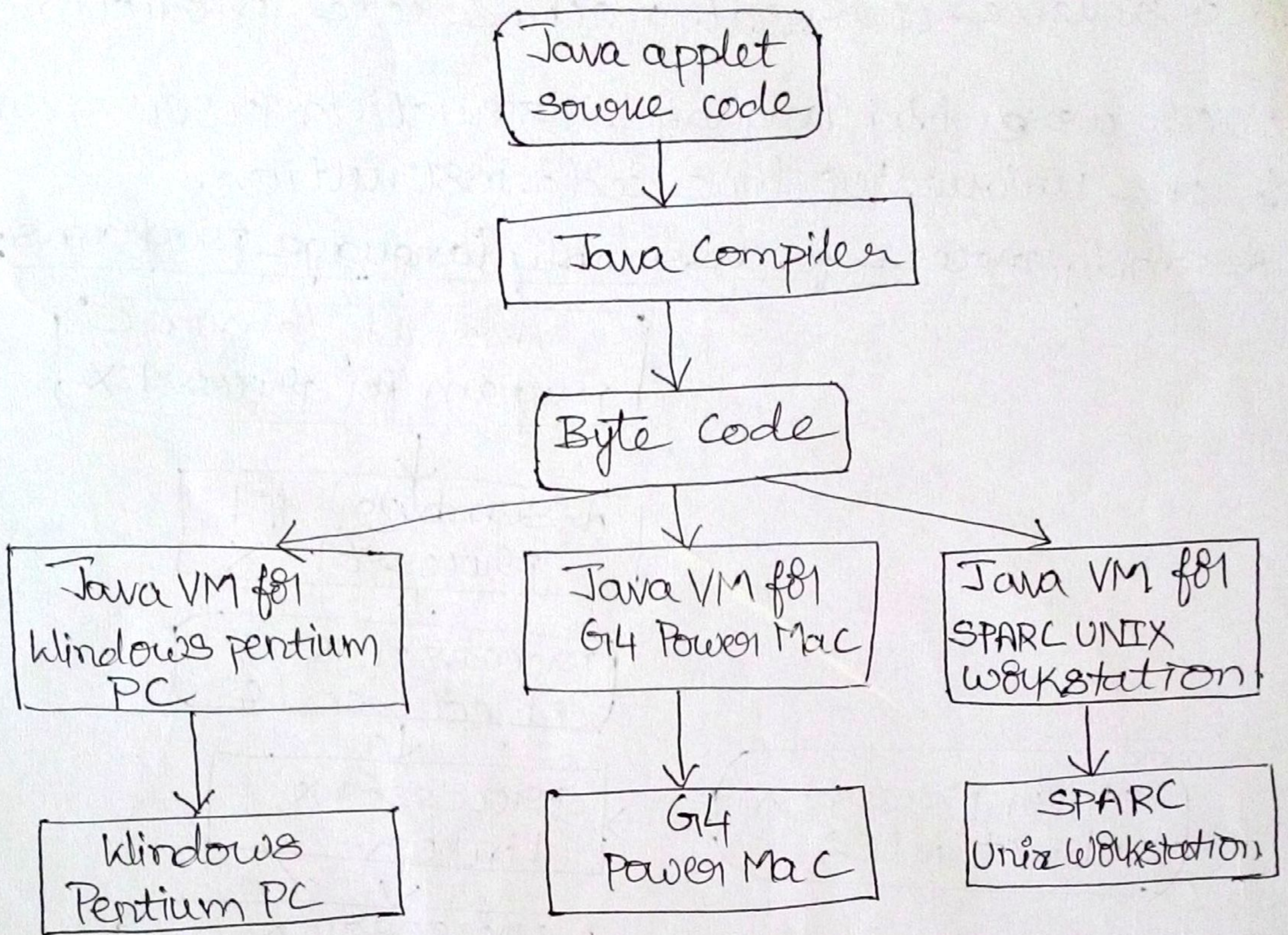
- A high-level language statement is usually converted to a sequence of several machine code instructions.
- Each assembly language instruction corresponds to one unique machine code instruction.

Assembly process for assembly language programs:-



- Many computers, however are much simpler than these systems. The microprocessor system used to control consumer appliances generally consist of a microprocessor and a small amount of memory and some circuitry to interact with a keypad and a digital display. In such a system, an assembly language program would be assembled to generate object code, but this might be the end of the process.

→ Compilation Process for Java applets :-



⇒ Assembly Language Instructions:-

- Assembly language instructions can be grouped together based on the types of operations they perform.

- Data Transfer Instructions
- Data operation Instructions
- Program control Instructions

→ Data Transfer Instructions.

- These instruction only copy the value to its destination.

- These instructions typically perform one of the following transfers.

- Load data from memory into the microprocessor.
- Store data from the microprocessor into memory.
- Move data within the microprocessor
- Input data to the microprocessor.
- Output data from the microprocessor.

→ Data Operation Instructions:-

- These perform some operation using one or two data values (operands) and store the result.

- Arithmetic instruction makeup a large part of the data operation instructions. (ADD, SUB, MUL, DIV).

- A special class of Arithmetic instructions are floating point instructions.

- Logic instructions perform basic logical operations on data. They are AND, OR or XOR.

- Shift instructions, shift the bits of a data value.

→ Program Control Instructions:-

- A Jump or branch instruction is commonly used to go to another part of the program.
- Assembly language instruction set may include instruction to call and return from subroutines.
- An assembly language instruction set may include specific instructions to generate interrupts; these are called software interrupts.
- Halt instruction causes a microprocessor to stop executing instructions, such as at the end of a program.

Data types:-

signed	-2^{n-1} to $2^{n-1} - 1$
unsigned	0 to $2^n - 1$

- float

- Boolean values TRUE & FALSE.

- character data.

characters are stored as binary values encoded using ASCII, EBCDIC, UNICODE.

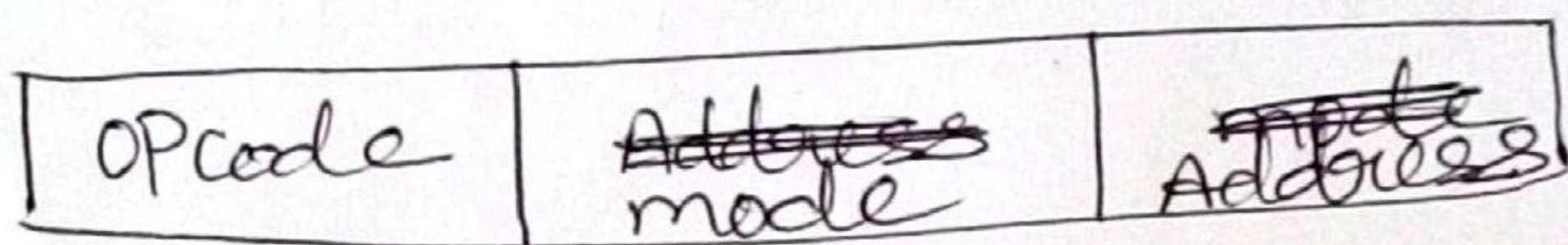
↓
American Standard Code for Information Interchange

EBCDIC - Extended Binary Coded Decimal Interchange code

UNICODE.

⇒ Instruction Formats :-

→ The control unit of CPU interpret each instruction code and provide the necessary control functions needed to process the instruction.



Instruction Format.

- Operation code (OPcode) - specifies the operation to be performed.
- Address - designates a memory address or a processor register.
- mode - designates the way the operand or the effective address is determined.

→ Computers may have instructions of several different lengths containing varying number of addresses. The number of address fields in the instruction format of a computer depends on the internal organization of its registers.

Most computers fall into one of three types of CPU organizations:

- ① Single accumulator organization
- ② General register organization
- ③ Stack organization.

① ⇒ uses one address field.

Eq:- ADD X

ADD - Addition

X - address of operand.

$AC \leftarrow AC + M[X]$.

② ⇒ needs three register address fields.

ADD R1, R2, R3

$R1 \leftarrow R2 + R3$

- if the destination register is the same as one of the source registers; the no: of address fields in the instruction can be reduced from three to two.

ADD R1, R2

$R1 \leftarrow R1 + R2$.

eg:

MOV R1, R2

denotes transfer $R1 \leftarrow R2$

③ ⇒ PUSH and POP instruction.

PUSH X

will push the word at address X to the top of the stack.

operation-type instructions do not need an address field in stack-organized computer.

ADD

this operation pop the top two numbers from stack, add, and push the sum onto the stack.

- To illustrate the influence of the number of addresses on computer programs, let us consider the evaluation of arithmetic statement

$$X = (A+B) * (C+D)$$

Three-Address Instructions:

ADD R1, A, B

$R1 \leftarrow M[A] + M[B]$

ADD R2, C, D

$R2 \leftarrow M[C] + M[D]$

MUL X, R1, R2

$M[X] \leftarrow R1 * R2$

Two-Address Instructions

MOV R1, A	$R1 \leftarrow M[A]$
ADD R1, B	$R1 \leftarrow R1 + M[B]$
MOV R2, C	$R2 \leftarrow M[C]$
ADD R2, D	$R2 \leftarrow R2 + M[D]$
MUL R1, R2	$R1 \leftarrow R1 * R2$
MOV X, R1	$M[X] \leftarrow R1$

One-Address Instructions :-

- use an implied accumulator (AC) register

LOAD A	$AC \leftarrow M[A]$
ADD B	$AC \leftarrow AC + M[B]$
STORE T	$M[T] \leftarrow AC$
LOAD C	$AC \leftarrow M[C]$
ADD D	$AC \leftarrow AC + M[D]$
MUL T	$AC \leftarrow AC * M[T]$
STORE X	$M[X] \leftarrow AC$

Zero-Address Instructions :-

TOS = TOP OF STACK.

PUSH A	$TOS \leftarrow A$
PUSH B	$TOS \leftarrow B$
ADD	$TOS \leftarrow (A+B)$
PUSH C	$TOS \leftarrow C$
PUSH D	$TOS \leftarrow D$
ADD	$TOS \leftarrow (C+D)$
MUL	$TOS \leftarrow (C+D) * (A+B)$
POP X	$M[X] \leftarrow TOS$

To evaluate arithmetic expressions in a stack computer, it is necessary to convert the expression into reverse Polish notation.

⇒ Eq:-

Evaluate the arithmetic statement:

$$X = \frac{A - B + C * (D * E - F)}{G + H * K}$$

- (a) Using a general register computer with three address instructions
- (b) with two address instructions
- (c) one address instruction, - accumulator type computer
- (d) using a stack organized computer with zero-address operation instructions